A Runtime Environment for Contract Automata



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Behavioural Contracts

- Behavioural types: model the behaviour of ensembles of services in terms of their interactions;
 - behavioural contracts, session types.
- Useful for:
 - reasoning formally about well-behaving properties,
 - building applications that are verified by construction against these properties.
- Contract automata: behavioural contracts modelled as FSA,
- In contract automata services match their requests and offers between each other to reach an agreement

Motivation

- few studies on how to derive the finalised software from the verified specification as behavioural contracts;
- these behavioural specifications are not yet a feature of standard mainstream programming languages.
- We investigate the connection between a behavioural specification given as contract automata and its implementation.

Contract Automata: Software Support

- CATLib : library implementing contract automata and their operations, for *specifying* applications using contract automata;
- CAT_App : Graphical front-end for designing contract automata;
- <u>CARE</u> : Runtime Environment for *implementing* applications specified via contract automata

https://github.com/contractautomataproject

CARE methodology



Operations of Contract Automata

- contract automata are composed to model all possible behaviours,
 - composition in agreement: all requests of contracts are matched by corresponding offers,
- if the composition does not satisfy the requirements or it is not in agreement, it is possible to automatically compute (i.e., **synthesise**) a refinement of the composition such that the requirements are satisfied,
- this refined automaton is called orchestration,
- the orchestration automaton implicitly assumes the presence of an orchestrator coordinating the services to execute the actions prescribed by the orchestration

Benefits of CARE

- We show a possible realisation of an orchestration engine,
 - abstracted away in the contract automata theory,
 - but needed for implementing applications specified with contract automata,
- improving our understanding of the relation between a specification with contract automata and its implementation, and the corresponding level of abstraction.
- **Benefits:** formal guarantees, reduction of the software complexity, separation of concerns, modularity.

Modal Service Contract Automata (MSCA)

- MSCA are FSA enhanced with:
 - Partitioned alphabet of actions:
 - offers !a (or \overline{a}) (A^{o}) and requests ?a (or a) (A^{r})
 - special idle action (• $\notin A^o \cup A^r$)
 - rank : the number of services in the contract,
 - Transitions partitioned into permitted (T^{\diamond}) and necessary (T^{\Box}) ,
 - Labels are list of actions and are constrained to be:



Contract Automata Runtime Environment

Contract Automata Runtime Environment

- the software is organised into:
 - the classes for the orchestrator and
 - the classes for the orchestrated services;
- offers and requests of contracts are an abstraction of low-level messages sent between the services and the orchestrator to realise them.

Runnable Orchestration, Runnable Orchestrated

- RunnableOrchestration:
 - abstract class implementing an orchestrator,
 - it reads the synthesised orchestration and orchestrates the RunnableOrchestratedContract to realise the overall application.
- RunnableOrchestratedContract:
 - abstract class implementing an executable wrapper responsible for the *composition of the specification of a service with its implementation*.
 - It implements a service that is always listening and spawns a parallel process when entering an orchestration.
 - During an orchestration, it receives action commands from the orchestrator or from other services, and it invokes the corresponding action method.

Two aspects to implement: action and choice



D.B. et al. (ISTI CNR Pisa)

Contract Automata Runtime Environment

CARE - match implementation



bob

alice



ror

Formal Guarantees : Adherence to the Contract

Algorithm 1 Orchestration	Algorithm 2 Service	Thread	
Require: non-empty orchestration automaton	Require: connected to the orchestrator		
Ensure: no exception is thrown	init Socket	\triangleright set socket timeout	
init Sockets > connect to the services	$cs \leftarrow initialState$	⊳ current state	
$cs \leftarrow initialState > current state$	while true do		
while true do	$act \leftarrow receive(socket)$		
$fws \leftarrow forwardStar(cs)$	if stop(act) then		
if empty(fws) & notFinal(cs) then	if final(cs) then		
throw Exception	return		
end if	else throw ContractViolationException		
$choice \leftarrow choice() \triangleright interact with services$	end if		
if $choice == stop \& final(cs)$ then	end if		
return	if choice(act) then		
end if	performChoice()	\triangleright interact with or-	
$tr \leftarrow select(fws, choice)$	continue	▷ chestration	
if tr not in agreement then	end if		
throw Exception	$tr \leftarrow select(forwardStar(cs), act)$		
end if	if no valid action then		
$doAction(tr)$ \triangleright interact with services	throw ContractViolationException		
$cs \leftarrow targetState(tr)$	else		
end while	invokeMethod(tr)		
	end if		
	$cs \leftarrow targetState(tr)$		
	end while		

if the orchestration is correctly synthesised no contract exception will ever be raised

Formal Guarantees : Uppaal model

- formal model covering aspects that are abstracted in the algorithms:
 - $\bullet\ centralised/distributed$ action
 - majoritarian/dictatorial choice
- TCP/IP sockets modelled through blocking FIFO buffers
- Difference between contract automata and Uppaal models:
 - *contract automata* are high level models of service applications that can be implemented using CARE,
 - CARE is a distributed application, suitable to be verified with model checking,
 - CARE itself is a software that has been formally modelled and verified using Uppaal.

Formal Guarantees : Uppaal model



Formal Model : Traceability



https://github.com/contractautomataproject/CARE/tree/master/src/spec/uppaal

Model Checking

• A[](not deadlock || (ror.Terminated && (forall(i:id_t) ROC(i).Terminated)))

- absence of deadlocks
- A[]((ror.Terminated && (forall (i:id_t) ROC(i).Terminated)) imply
 allEmpty())
 - absence of orphan messages
- A[](ror.Stop imply A<>(ror.Terminated && (forall (i:id_t) ROC(i).Terminated) && allEmpty()))
 - whenever a choice to stop is made, eventually all services and the orchestrator will terminate their execution.

CARE evaluation

- two examples: Alice and Bob, Composition service
- two implementations: one with CARE, one without using CARE

		LOC	Cyclomatic Complexity	Cognitive Complexity
Alice and Bob	without CARE	777	134	166
	with CARE	153	16	8
Composition Service	without CARE	854	155	211
	with CARE	279	42	55

Conclusion, Future work

- We have presented the first runtime environment for contract automata, called CARE.
- Future work support for choreographies and product lines is formalised in contract automata and implemented in CATLib, but not supported by CARE
- Model-based testing
 - Uppaal tests generation,
 - generation of tests at the level of contract automata.
- Other enhancements to do

https://github.com/orgs/contractautomataproject/projects/1

Other references

- Basile, D. and ter Beek, M.H., 2022. Contract automata library. Science of Computer Programming (Original Software Publication).
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